## Relative Monadicity

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## Part I

#### Categories of structures

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What does it mean for D to be a category of Structured objects in E?

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At minimum, we expect U to be faithful, but typically we are interested in something stronger.

## Algebraic structure

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Categories of algebras for monads are very well understood, and so it is frequently useful to exhibit categories as monadic over well-behaved categories.

Let T be a monad on a category  $\mathcal{E}$ . A T-algebra is an object  $E \in \mathcal{E}$  equipped with, for each

$$X \longrightarrow E$$

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$$\begin{array}{c} f \\ \uparrow \\ \uparrow \\ \chi \\ \downarrow \\ \uparrow \\ \xi \end{array}$$

satisfying two laws.

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f given an interpretation of X in E

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satisfying two laws.

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A T-algebra is an object  $E \in E$  equipped

with, for each

TX, 
$$f^{\varepsilon}$$
 we can interpret all these new terms in E

X  $\longrightarrow$  E

satisfying two laws.

## Arity

Because we must be able to extend morphisms  $X \xrightarrow{f} E$  with arbitrary domain, E must potentially interpret operations with very large (e.g. infinitary) arity.

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However, in many situations we are only concerned with interpreting terms with small (e.g. finite) arity, for which the algebras for a monad are too restrictive.

#### Relative monads

A <u>relative monad</u> is a generalisation of a monad that captures the intuition that the operations may have restricted arity.

For  $A \xrightarrow{J} \mathcal{E}$ , a J-relative monad comprises: • a functor  $A \xrightarrow{T} \mathcal{E}$ • a natural transformation  $J \Rightarrow T$ • for each  $J \times f \to T \times f$ • for each  $J \times f \to T \times f$ • a  $T \times f \to T \times f$ 

## Algebras for a relative monad

Given a  $A \xrightarrow{J} E$ -relative monad T, a T-algebra is an object  $E \in E$  equipped with, for each

$$\begin{array}{c} TX \\ 1x \\ TX \\ \downarrow \\ f \end{array}$$

satisfying two laws.

## Algebraic theories

A suggestive class of examples arise in universal algebra. Monads relative to FinSet \( \rightarrow Set \) are in bijection with algebraic theories.

Their algebras therefore only require extensions along functions  $X \longrightarrow E$  with X finite, corresponding to finitarity of operations in universal algebra.

## Categories of structures II

Given a category & and a functor  $A \xrightarrow{\mathcal{I}} \mathcal{E}$ , the categories of algebras for J-relative monads capture the intuition of categories whose objects are those of & equipped with operations with arity in A, subject to equations.

## Categories of structures II

Given a category  $\mathcal{E}$  and a functor  $\mathcal{A} \xrightarrow{\mathcal{I}} \mathcal{E}$ , the categories of algebras for J-relative monads capture the intuition of categories whose objects are those of & equipped with operations with arity in A, subject to equations.

How can we characterise them?

#### Monadicity theorems

A functor  $D \xrightarrow{\Sigma} \Sigma$  is monadic iff it admits a left adjoint and either:

Beck (c. 1966) U creates U-split coequalisers

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or:

Paré (1971) U creates absolute colimits

#### Absolute colimits

Given a functor  $X \xrightarrow{D} E$  admitting a colimit, colim D is absolute if either:

- · it is preserved by every functor from &
- it is preserved by the presheaf embedding  $\mathcal{E} \hookrightarrow [\mathcal{E}^{\circ p}, \mathsf{Set}]$

## Relative monadicity

Given functors

$$\frac{0}{J} \in \mathcal{E}$$

we would like conditions under which U exhibits D as being the category of algebras for a J-relative monad.

#### Relative adjunctions

Given functors

$$A \xrightarrow{\Gamma} \Sigma$$

We say L is <u>left J-adjoint</u> to R  $(L \to_J R)$  if  $\mathcal{C}(L \to \chi, y) \cong \mathcal{E}(J \to \chi, Ry)$ 

natural in x, y.

## J-absolute colimits

Given a functor  $A \xrightarrow{J} \Sigma$  and a functor  $X \xrightarrow{D} \Sigma$  admitting a colimit, colim D is  $\underline{J}$ -absolute if it is preserved by the restricted presheaf embedding

$$\mathcal{E} \longrightarrow [\mathcal{E}^{\circ P}, Set] \longrightarrow [\mathcal{A}^{\circ P}, Set]$$

#### The relative monadicity theorem

Consider functors:

$$\frac{1}{\sqrt{1}} \in \mathcal{E}$$

where J is dense.

Theorem U is J-monadic iff it admits a left J-adjoint and creates J-absolute colimits.

## Applications

Various frameworks for <u>monad-theory</u> correspondences have been proposed in recent years, the most popular being:

- · Monads with arities
- · Nervous monads

By the results of [Ark 22], these are all subsumed by relative monads.

## Monadicity with arity & nervous monadicity

Consequently, the relative monadicity theorem exhibits monadicity theorems for monads with arities and nervous monads.

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Corollary Let  $A \xrightarrow{J} \mathcal{E}$  and  $D \xrightarrow{U} \mathcal{E}$  be functors. U exhibits D as the category of algebras for a monad with arities J / J-nervous monad iff U admits a left adjoint and creates J-absolute colimits.

## Part I

# When is the composite of two functors monadic?

Suppose we are given two functors  $A \xrightarrow{U_1} B \xrightarrow{U_2} C$ 

and told they are both monadic.

We cannot conclude  $A \xrightarrow{U_2U_1} G$  is monadic (e.g.  $Cat \rightarrow Grph \rightarrow Set^2$ ).

# When is the composite of two functors monadic?

There are various <u>heuristics</u> for detecting monadicity of composites (e.g. crude monadicity), but these only provide sufficient conditions.

We would prefer a precise characterisation.

## The pasting law for pullbacks

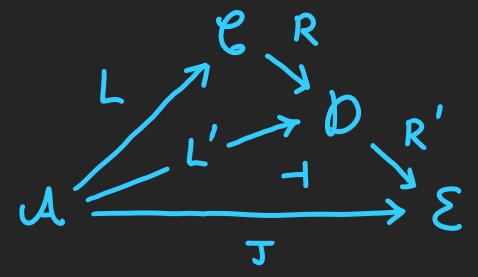
Suppose we are given a diagram of morphisms as follows.

$$\begin{array}{c} A \longrightarrow B \longrightarrow C \\ \downarrow & \downarrow & pb \downarrow \\ X \longrightarrow Y \longrightarrow Z \end{array}$$

Then the outer rectangle is a pullback iff the left square is a pullback.

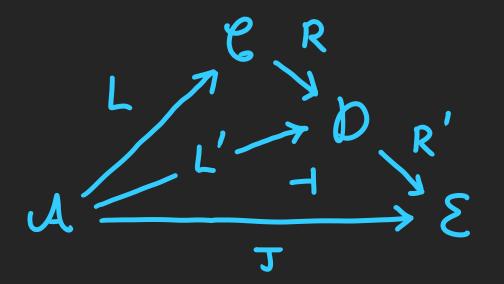
## The pasting law for relative adjunctions

Suppose we are given a diagram of functors as follows.



The outer triangle is a relative adjunction iff the left triangle is a relative adjunction.

## The pasting law for relatively monadic adjunctions



Theorem If R' is J-monadic, then R'R is J-monadic iff R is L'-monadic.

# When is the composite of two functors monadic?

Suppose we are given two functors:  $A \xrightarrow{U_1} B \xrightarrow{U_2} C$ 

Corollary If  $U_2$  is monadic, then  $U_2U_1$  is monadic iff  $U_1$  is  $F_2$ -monadic.

#### Useful corollaries

Corollary 1 Given adjunctions,

if U2 and U2U, are monadic, so is U1.

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if U2 and U2U1 are monadic, so is U1.

Corollary 2 For every morphism  $S \Rightarrow T$ between finitary monads on a LFP category the induced functor  $T-Alg \rightarrow S-Alg$  is monadic.

## Summary

- Relative monads permit us to capture structures interpreting operations of restricted arity.
- · We establish two monadicity theorems.
  - A Paré-style characterisation in terms of absolute colimits
  - A pasting law for relatively monadic adjunctions

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